# Fast Fourier Transform Reductions for Bayesian Network Inference Vincent Hsiao<sup>1</sup>, Dana Nau<sup>1</sup>, Rina Dechter<sup>2</sup>



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# Background

# **Bayesian Network (BN)**

- Bayesian Network: graphical model (X,D,F)
- $\circ$  Variables:  $X = \{X_1, X_2, ..., X_N\}$
- $\circ$  Domains:  $D = \{D_{X_1}, D_{X_2}, \dots, D_{X_N}\}$
- Parent Functions:  $F = \{F_1, F_2, \dots, F_N\}$

# Causal Independence (CI)

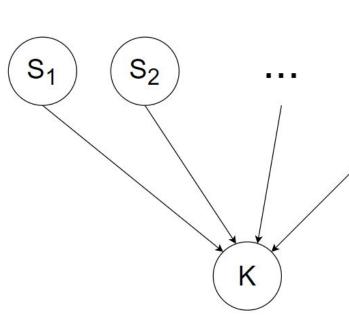
- Probabilistic relationship between a set of causes  $\{c_1,...,c_n\}$  and an effect e, such that:  $e = h_1 * h_2 ... * h_n$
- where each hidden variable h, is some probabilistic function of its corresponding c<sub>i</sub> and \* is some commutative and associative binary operator
- Summation-type CI: operator "\*" is addition "+"
- Network Transformation: Given a CI BN fragment {X,D,F} with a set of causes  $\{c_1,...,c_n\}$ , an effect e, and hidden variables  $\{h_1,\ldots,h_n\}$ , a network transformation is a new network {X',D',F'} constructed over some computational ordering:
- $e = (\dots(((h_1 * h_2) * h_3) * h_4) * \dots) * h_n$ with new intermediate variables:

$$y_i: \{y_1 = h_1 * h_2, y_2 = y_1 * h_3 \ldots \}$$

 Algorithms such as ci-elim-bel (Zhang and Poole, 1996; Rish and Dechter, 1998) exploit network transformations to accelerate bucket elimination

#### **Applications**

- Distributed Computing, Fault Tree Analysis
- N different resource providers with stochastic availability (k-out-of-n model)



k/n sum network

Evolutionary Game Theory (source-sink networks)

# **Problem Statement**

# **Current Approach**

 Network transformations (temporal decomposition), ci-elim-bel

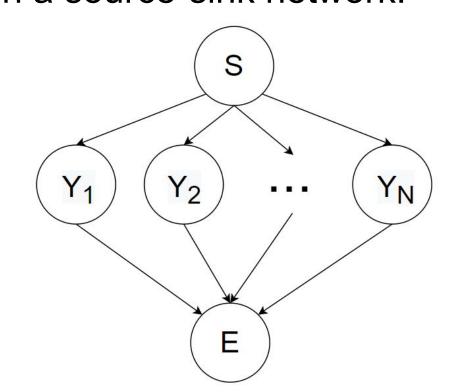
#### Proposed Approach

- Eliminate unneeded nodes in summation-type CI efficiently through the use of the FFT
- Integrate the FFT with bucket elimination into an efficient inference algorithm

# FFT Reduction

# **Computing Random Variable Sums**

Given a source-sink network:



source-sink network (2)

The size of the parent function for the node E expressed as a conditional probability table (CPT) is exponential in the number of Y nodes

We would like to reduce the size of the CPT of the network. This is useful in applications where one wants to find the marginal distribution of node E such as in:

- Test score prediction
- Evolutionary games

It can be shown that the distribution P(E|S) can be expressed as a convolution of individual distributions  $P(Y_i|S)$ 

From the convolution theorem and the application of the FFT, we know that:

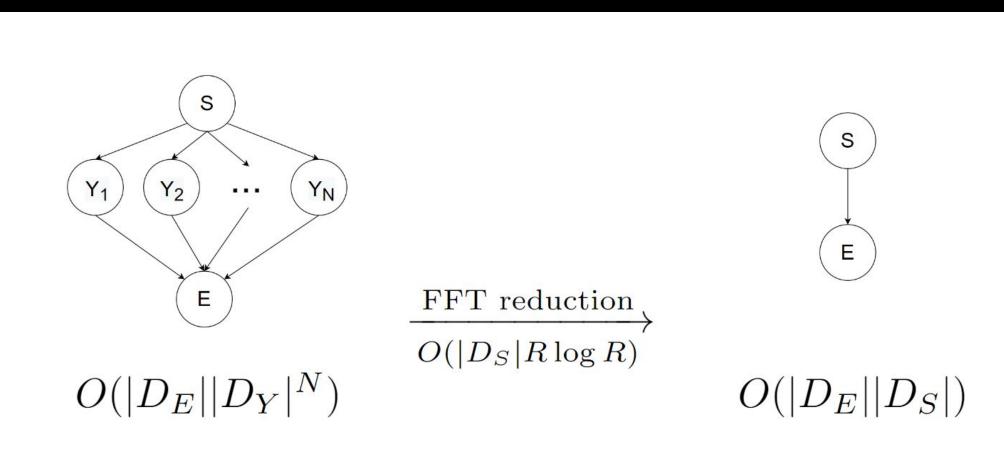
**Theorem 2.4** (FFT Time Complexity). For i.i.d random variables  $X = \{X_1, X_2, \dots, X_N\}$  where each variable has a domain size |D| and their sum  $Z = \sum X_i$ , computing P(Z = k) using the FFT will take time  $O(N|D|\log(N|D|))$ . For non-i.i.d variables, the time complexity is  $O(N^2|D|\log(N|D|))$ .

#### It follows that:

**Theorem 3.1** (FFT Reduction). Let  $B = \{X, D, F\}$ with  $X = \{S, Y_1, \dots, Y_N, E\}$  be a source-sink network with N i.i.d paths. The network can be transformed into  $\{X', D', F'\}$  such that  $X' = \{S, E\}$  reducing the CPT for E from size  $O(|D_E||D_Y|^N)$  to size  $O(|D_S||D_E|)$  in  $O(|D_S|R\log R)$  time where R is the numerical range of random variable E.

Furthermore, it can be shown that certain computations performed in the algorithm *ci-elim-bel* can also be formulated as convolutions, leading to an improved algorithm: ci-elim-FFT

# FFT Reduction



# **CI-Elim-FFT**

Algorithm 2: ci-elim-FFT Input: A Bayesian network B, evidence eOutput:  $P(x_1|e)$ 

Generate a decomposition network from BGenerate ordering  $o = \{Z_1, \ldots, Z_n\}$  with  $Z_1 = \{x_1\} \text{ using } B'$ for  $i = n \rightarrow 1$  do // create buckets

 $\forall x \in Z_i$ , put all network functions with x as highest ordered variable in  $bucket_i$ enc

for  $i = n \rightarrow 1$  do // process buckets  $// h_1, \ldots, h_m$  are functions in bucket<sub>i</sub> if  $Z_i = \{u_l; u_k\}, u = u_l + u_k$  then  $h^{u_l} = \prod_{j,u_l \in h_j} h_j \; ;$  $h^{u_k} = \prod_{j,u_k \in h_j} h_j$ ;  $h^{Z_i} = \mathcal{F}^{-1}\{\mathcal{F}\{h^{u_l}\}\cdot\mathcal{F}\{h^{u_k}\}\}$ 

use regular ci-elim-bel to compute  $h^{Z_i}$ Put  $h^{Z_i}$  in the highest bucket that mentions  $h^{Z_i}$ 's variable.

Return  $\alpha h^{x_1}$ , ( $\alpha$  is a normalizing constant).

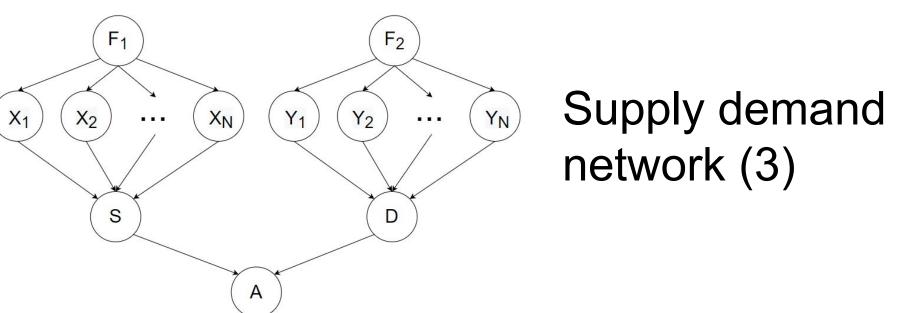
# Complexity on k/n sum network (1)

	$i.i.d Y_i's$	non- $i.i.d Y_i's$
BE	$O(N D_X R_Z^2)$	$O(N D_X R_Z^2)$
FFT	$O( D_X R_Z\log R_Z)$	$O(N D_X R_Z\log R_Z)$

# **Experimental Setup**

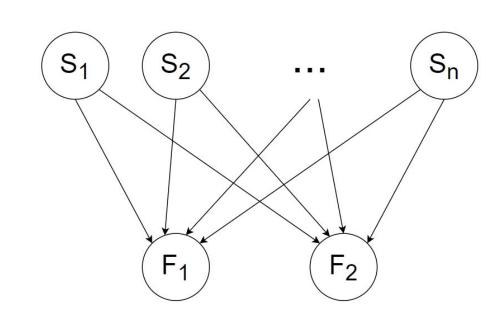
# **Experimental Setup (Set 1)**

- Evaluate FFT reduction for inference on a selection of networks with summation-CI
- Compare with:
- Vanilla Bucket Elimination (Naive)
- Temporal Decomposition (Temporal)
- Test inference on three types of networks (1), (2), (3)



# **Experimental Setup (Set 2)**

• Evaluate *ci-elim-FFT* compared to *ci-elim-bel* on general two layer additive networks (4):

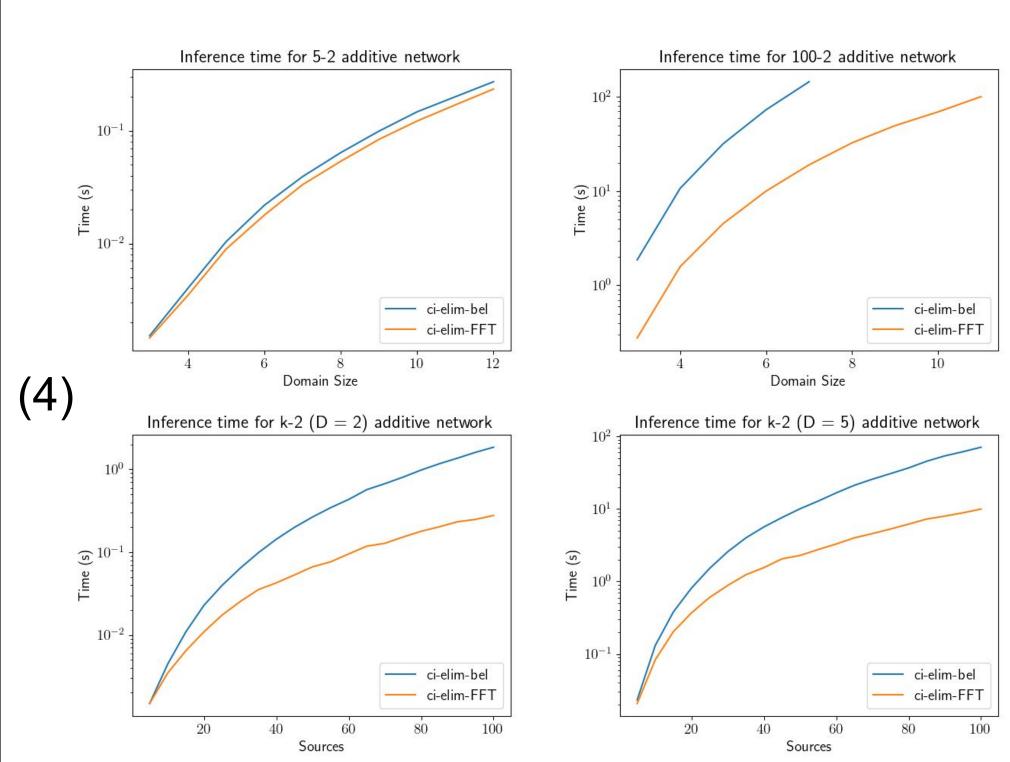


Two layer additive network (4)

# Results

The FFT reduction method obtains a significant computational advantage over the naive approach as well as inference on the temporal decomposition

(3)



ci-elim-FFT demonstrates better scaling with respect to number of source nodes and domain size in general additive networks compared to *ci-elim-bel* 

# **Future Research**

 Explore integration into existing lifted variable elimination algorithms

# Acknowledgements

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