

AND/OR Branch-and-Bound for Computational Protein Design Optimizing K*

Bobak Pezeshki (pezeshkb@uci.edu), Radu Marinescu (radu.marinescu@ie.ibm.com), Alex Ihler (ihler@ics.uci.edu), Rina Dechter (dechter@ics.uci.edu)

Summary

We introduce AOBB-K*MAP, a new branch-and-bound algorithm over AND/OR search spaces for solving the K*MAP protein design problem. In addition to formulating CPD as a graphical model for K* optimization and providing a new efficient algorithm, we also introduce a statically compiled heuristic for K*MAP not previously used in CPD. This work extends algorithms for Marginal MAP (MMAP) and provides a framework and proof-of-concept for continued adaptation of existing state-of-the-art mixed inference schemes over AND/OR search spaces to protein design.

$$\mathcal{M} = \langle X, D, F \rangle$$

Formulation

$$X = R \cup C_\gamma$$

Two sets of variables:

- Residue variables – capture the amino acid (aa) at the respective position in the protein
- Conformation variables - side chain conformation of the aa at the corresponding residue for protein states $\gamma \in \{\text{Bound}, \text{Unbound}\}$

$$D = D_R \cup D_{C_\gamma}$$

Domains:

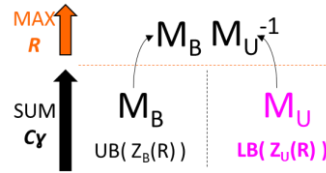
- D_R - amino acids being considered at each position
- D_{C_γ} - various side chain conformations for all aa considered at the corresponding residue for each protein state γ

$$F = E_\gamma \cup C$$

Functions:

- Energy of interactions between the amino acids for each protein state γ
- Constraints enforcing consistency between side chain conformations and amino acid assignments

MBE-K*



Based on a constrained ordering such that conformation variables are processed first, messages corresponding to upper and lower bounds on the partition functions of the conformation variables associated with the various protein states are computed and passed up to be maximized over

Algorithm 2: AOBB-K*MAP

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input : CPD graphical model  $\mathcal{M}$ ; pseudo-tree  $\mathcal{T}$ ;
        $K^*$  upper-bounding heuristic function
        $h_{K^*}^{ub}(\cdot)$ ;  $Z_\gamma$  upper-bounding heuristic
       function  $h_{Z_\gamma}^{ub}(\cdot)$ ; and subunit stability
       threshold  $threshold(\gamma)$  for each subunit
        $\gamma \in \varphi$ 
output:  $K^*MAP(\mathcal{M})$ 
1 begin
2   Initialize MiniSat with constraints from  $\mathcal{M}$ 
3   and generate literals via constraint propagation
4    $\pi \leftarrow$  dummy AND node  $n_D$ 
5    $ub_{K^*}(n_D) \leftarrow \prod_{m \in ch_{\mathcal{T}}(n_D)} h_{K^*}(m)$ 
6    $lb_{K^*}(n_D) \leftarrow -inf$ 
7    $g(n_D) \leftarrow 1$ 
8   foreach  $\gamma \in \varphi$  do
9      $A_{Z_\gamma}^x(n_D) \leftarrow 1$ 
10     $A_{Z_\gamma}^z(n_D) \leftarrow 0$ 
11     $ub_{Z_\gamma}(n_D) \leftarrow \prod_{m \in ch_{\mathcal{T}_\gamma}(n_D)} h_{Z_\gamma}(m)$ 

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while  $n_X \leftarrow EXPAND(\pi)$  do
  if MiniSat( $\pi$ ) = false then
    | PRUNE( $\pi$ )
  else if  $\exists \gamma \in \varphi$  s.t.
    |  $UB_{Z_\gamma}(n_X) < threshold(\gamma)$  then
    | | PRUNE( $\pi$ )
  else if  $X \in R$  then
    | if  $\exists a \in anc^{OR}(n)$  s.t.
    | |  $ub_{K^*}(a, \pi) < lb_{K^*}(a)$  then
    | | | PRUNE( $\pi$ )
    | end
  else if  $ch_{\mathcal{T}}^{unexp}(n) = \emptyset$  then
    | | BACKTRACK( $\pi$ )
  end
return
   $ub_{K^*}(n_D) = lb_{K^*}(n_D) = K^*MAP(\mathcal{M})$ 
end

```

Results and Future Work

Problem	iB	X	Dmax	w*	d	UB	OR	AND	CPP	UBP	SSP	EH	time	K*MAP	BBK* t	BBK* sln
1gwc_00021	4	12	203	4	6	10.29	28766	134930	77823	55	2	5	16	9.79	152	9.79
2hnu_00026	4	14	203	5	7	15.08	22010	105458	76657	38	0	4	7	13.18	437	13.18
2hmv_00025	4	16	203	6	8	15.04	115194	297138	84882	39	0	3	16	13.65	962	13.65
2rf9_00018	6	18	205	7	9	16.68	20137	85033	87306	78	0	4	15	15.79	187	15.79
2rfd_00035	6	16	205	6	8	17.70	896239	4253159	3273123	40	0	4	381	17	1242	16.77
2rfe_00030	4	14	203	5	7	11.53	20393	164126	359007	87	40	7	19	10.50	182	10.50
2rfe_00043	6	16	203	6	8	18.48	15390	40297	422357	34	43	4	80	18.04	50	18.04
2rfe_00044	6	16	203	6	8	18.62	37887	99927	1047107	30	3	6	86	18.19	75	18.19
2rl0_00008	4	10	203	3	5	11.16	2	3	0	40	0	3	3	11.16	262	9.46
2xgy_00020	4	14	203	5	7	11.47	43643	262523	743860	40	0	2	14	10.60	887	10.60
3cal_00032	6	16	203	6	8	13.38	133851	1067419	531976	32	6	4	125	11.62	1429	11.62
3u7y_00009	5	12	203	4	6	4.51	2	3	0	40	0	3	6	4.51	191	4.51
4kt6_00023	4	16	203	6	8	14.80	38186	101546	23877	16	19	4	7	12.69	136	12.69
4wwi_00019	5	14	203	5	7	15.43	8094	30774	17888	40	0	2	7	14.99	26	14.99
1gwc_00021	4	13	203	4	7	12.51	33881	590621	473189	388	6	8	205	11.92	551	11.72
2hmv_00025	4	17	203	6	9	18.38	215171	550559	220825	77	0	4	153	16.18	880	13.65
2rfe_00012	5	15	205	5	8	14.36	3127	10003	32610	57	0	3	85	13.93	12	13.93
2rfe_00014	5	15	205	5	8	14.79	4087	13087	39411	57	0	3	85	14.36	45	14.36
2rfe_00017	5	15	203	5	8	11.46	245894	1063198	6389737	227	25	43	333	10.86	78	10.80
2rfe_00030	4	15	203	5	8	13.61	256957	1327425	2816050	726	83	77	274	11.12	275	10.97
2xgy_00020	5	15	203	5	8	11.39	398102	2383318	7422285	42	0	20	360	10.96	1388	10.96
3u7y_00009	4	13	203	4	7	4.95	36760	228568	564654	204	7	28	99	4.51	216	4.51
3u7y_00011	4	13	203	4	7	12.29	5758	16108	68579	50	0	5	86	11.85	27	11.85
4wwi_00019	5	15	203	5	8	16.05	22945	87485	91677	176	75	5	180	14.99	34	14.99

Preliminary Results:

- AOBB-K*MAP performs better than BBK* time-wise for the majority of problems tested with 2 MAP variables and better than half of those with 3 MAP variables
- wMBE-K* provides solutions to conditioned subproblems during search, sometimes able to solve the entire K*MAP problem exactly
- AOBB-K*MAP finds solutions with greater K* values than BBK*

Future Work:

- Incorporation of constraints into wMBE-K* to exploit determinism, and the use of a dynamic heuristic
- Adaptation of state-of-the-art [approximate] search and sampling algorithms to solving the K*MAP problem.

K* Objective

$$K^*(R_1, \dots, R_N) = Z_B(R_1, \dots, R_N) / Z_U(R_1, \dots, R_N)$$

$$Z_\gamma(R_1 \dots R_N) = \sum_{C_1, \dots, C_N} \prod_{\mathcal{C}_{\gamma(i)} \in \mathcal{C}} \mathcal{E}_{\gamma(i)}(R_i, C_{\gamma(i)}) \cdot \prod_{E_{\gamma(i)}^{sb} \in E^{sb}} e^{-\frac{E_{\gamma(i)}^{sb}(C_{\gamma(i)})}{\mathcal{Z}_T}} \cdot \prod_{E_{\gamma(ij)}^{pw} \in E^{pw}} e^{-\frac{E_{\gamma(ij)}^{pw}(C_{\gamma(i)}, C_{\gamma(j)})}{\mathcal{Z}_T}}$$

$$\text{task: } K^*MAP = \max_{R_1, \dots, R_N} K^*(R_1, \dots, R_N)$$

